

## theory of Automata Lecture Notes: 1. Introduction & Theory of Formal Language

Computers run the world.

I am the cheese.

I never tell lies.

Now consider the derivation (construction) of the first sentence using the above grammar.

sentence		subject	verb-phrase	object
	F <sub>0</sub>			
	D <sub>1</sub> E			
	F <sub>0</sub>	This	verb-phrase	object
	D <sub>1</sub> E			
	F <sub>0</sub>	This	verb	object
	D <sub>1</sub> E			
	F <sub>0</sub>	This	is	object
	D <sub>1</sub> E			
	F <sub>0</sub>	This	is	a noun
	D <sub>1</sub> E			
	F <sub>0</sub>	This	is	a university
	D <sub>1</sub> E			

In addition to several reasonable sentences, some can also derive nonsense sentences like 'Computers run cheese' or 'This am a lies'. These sentences don't make semantic sense, but they are syntactically correct because they are of the sequence of **subject**, **verb-phrase**, **verb** and **object**.

It is very difficult to define the complete language with a finite number of rules. It is difficult to list all acceptable sentences of a language. In general the language should have the following properties:

- Well defined, without ambiguity.
- Using formula, we should be able to recognize in a finite time, whether any given word is in the language or not.

### Formal Definition of Grammar:

A formal definition of a grammar G can be given in 4-tuples as:

$G = (N, \Sigma, P, S)$  where

N is a finite set of non-terminals;

$\Sigma$  or T is a finite nonempty set of terminals;

S is the start symbol and  $S \in N$

P is a finite set of productions of the form:  $\alpha \rightarrow \beta$  as given in the above English grammar.

Definitions of the terms used above & some other grammar related terms are described as:

Symbols: A symbol means a point, letters, digits etc.

Alphabet: An alphabet ( $\Sigma$ ) is a finite, nonempty set of fundamental units, set of letters, character or symbols (Cohen pp-8)  
i.e.  $\Sigma = \{a, b, c, \dots, z\}$  i.e. an alphabet of cardinality 26.